Greedy Algorithms

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Overview

- change-making problem: give change for a specific amount n with the least number of coins of the denominations $d_1 > ... > d_m$
 - greedy approach: use largest denomination that is less than *n*, and reduce remainder. Repeat until the total of the change is equal to the change required.
- **greedy technique**: suggests constructing solution through sequence of steps, each expanding partial solution, until a complete solution is reached
 - for each step, the choice made must be:
 - * feasible, i.e. satisfies problem constraints
 - * locally optimal
 - * irrevocable: cannot be changed on subsequent steps
- in some instances greedy technique does provide globally optimum solution
 - minimum spanning tree problem: Prim's, Kruskal's algorithms
- often greedy technique will not yield optimum solution, but may still be useful to produce approximate solution
- Dijkstra's algorithm: shortest-path in a weighted graph
- · Huffman trees: data compression method
- greedy algorithms are typically intuitive and simple
- often difficult to prove they are optimal (if they are), approaches include:
 - mathematical induction, or
 - show that on each step greedy approach does at least as well as any other algorithm could
 - show optimal solution (rather than optimal efficiency)

Minimum Spanning Tree Problem

- **spanning tree**: connected acyclic subgraph (i.e. tree) containing all vertices of an undirected connected graph
- minimum spanning tree: spanning tree of smallest weight for a weighted graph
- weight: sum of weights on all edges
- given *n* points, connect them in the cheapest possible way, such that there is a path between every pair of points
- arises in all sorts of networks: communications, computer, transportation, electrical
 - cheapest way to achieve connectivity
 - identifies clusters in data sets
 - classification in archaeology, biology, ...
 - helpful for constructing approximate solutions to more difficult problems e.g. travelling salesman problem
- exhaustive search approach
 - exponential growth with graph size (especially for dense graphs)
 - generating all spanning trees is more difficult than finding minimum spanning tree

Prim's Algorithm

- construct minimum spanning tree through sequence of expanding subtrees
- initial subtree consists of arbitrary vertex
- algorithm expands the tree greedily by attaching nearest neighbour not in the tree
 - nearest: vertex out of tree connected to vertex in tree by edge of smallest weight
- algorithm stops when all vertices are in the tree
- total iterations: n-1

```
1 Prim(G)
2 # Prim's algorithm for constructing minimum spanning tree
3 # Input: weight graph G=<V,E>
4 # Output: E_T, set of edges composing minimum spanning tree of G
5 V_T = {v_0} # arbitrary vertex
6 E_T = empty_set
7 for i from 1 to |V|-1:
8 find minimum weight edge e* = (v*, u*) among all edges (v, u), such
8 that v is in V_T,
9 and u is in V-V_T
10 V_T = union(V_T, u*)
```

```
11 E_T = union(E_T, e*)
12 return E_T
```

- implementation requires each vertex not in the current tree to have information about shortest edge connecting it to vertex: attach labels to vertex
 - name of nearest vertex
 - weight (length) edge
- if no adjacent vertices:
 - name: ∞
 - label: null

Dijkstra's Algorithm